



# A Language Stack for implementing Contracts

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Context

# SMART

**An actual  
contract,  
executed  
automatically.**

# CONTRACT

**Any**  
Turing Complete  
**Program**  
**running on**  
**a Blockchain.**

over time



# **An actual contract, executed automatically.**

**Multiple Parties.**

**Decision ||**

**Agreement ||**

**Coordination.**

**(Legally) Binding &  
Trusted.**

**Formal Language.**

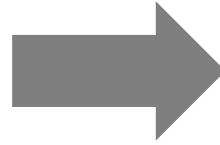
**Checkable.**

**Understandable.**

**„Event Tracking“**

**Progress over time**

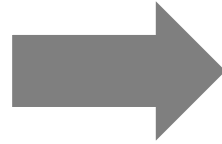
**Contract  
Definition**



**Contract  
Execution**



**Contract  
Definition**

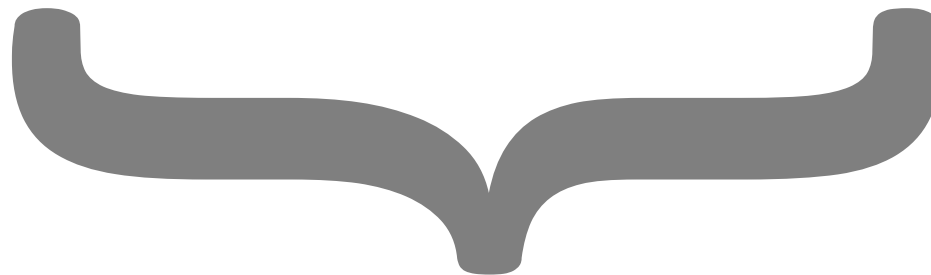


**Contract  
Execution**

Understand Behavior  
Functional Correctness

Non-Repudiability  
Verified Behavior  
Non-Gameability

**BC**



**TRUST**

# Blockchains

can provide certain  
**non-functional properties**  
to executable contracts.

---

# Blockchains

are a suitable (partial)  
**implementation technology**  
iff these properties are needed.

# Verification

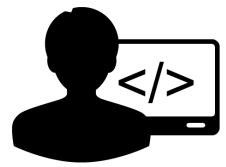
Ensure that the program performs correctly the things the program text tells it to do.

# Validation

Ensure that the program does the correct things, wrt. to the requirements.



# Verification



**Contract  
Execution**

Ensure that the program performs correctly the things the program text tells it to do.

# Validation



**Contract  
Definition**

Ensure that the program does the correct things, wrt. to the requirements.



# Correct-by-Construction

The language/framework/  
API/modeling tool doesn't allow  
a particular class of mistakes.

## Analysis-and-Fix

You analyze the code/model after the  
fact and try to find problems which  
devs then fix.

# **Correct-by-Construction**

## **Languages**

## **Analysis-and-Fix**

## **Analysis Tools**

Formal Language.  
Checkable.  
Understandable.

# DSL

Domain  
Specific  
Language







**Languages**

**Analysis Tools**



# Languages

## Analysis Tools



State machines can **always** be checked for dead states and unused transitions.

Decision tables must **always** be overlap-free and complete.





# Lots of History & Research



## **Computational Law**

Obligation, Permission

Ordering, Causality, Time

Event, State

# Lots of History & Research



## **Composing contracts: an adventure in financial engineering**

<https://lexifi.com/files/resources/MLFiPaper.pdf>



## **POETS Process-oriented event-driven transaction systems**

<https://github.com/legalese/poets/blob/master/doc/Henglein%20-%20POETS%20Process-oriented%20event-driven%20transaction%20systems.pdf>



## **Contracts in Programming and in Enterprise Systems**

<https://github.com/legalese/poets/blob/master/doc/hvitvedmaster.pdf>



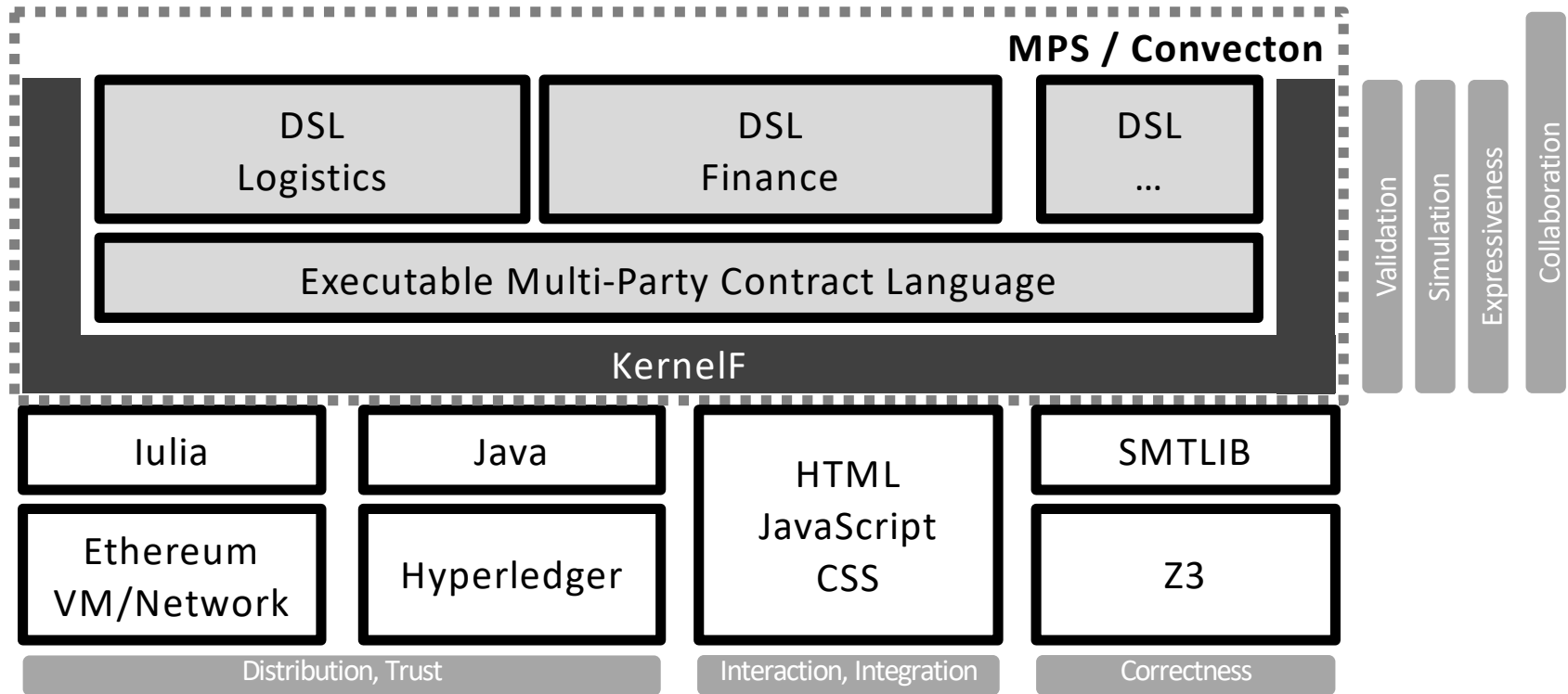
## **Domain-Specific Languages for Enterprise Systems**

<https://bitbucket.org/jespera/poets/raw/c0ee7194ce57d2ad6ca88948a44e88e546d5f4a/doc/poets-techreport/tr.pdf>

2

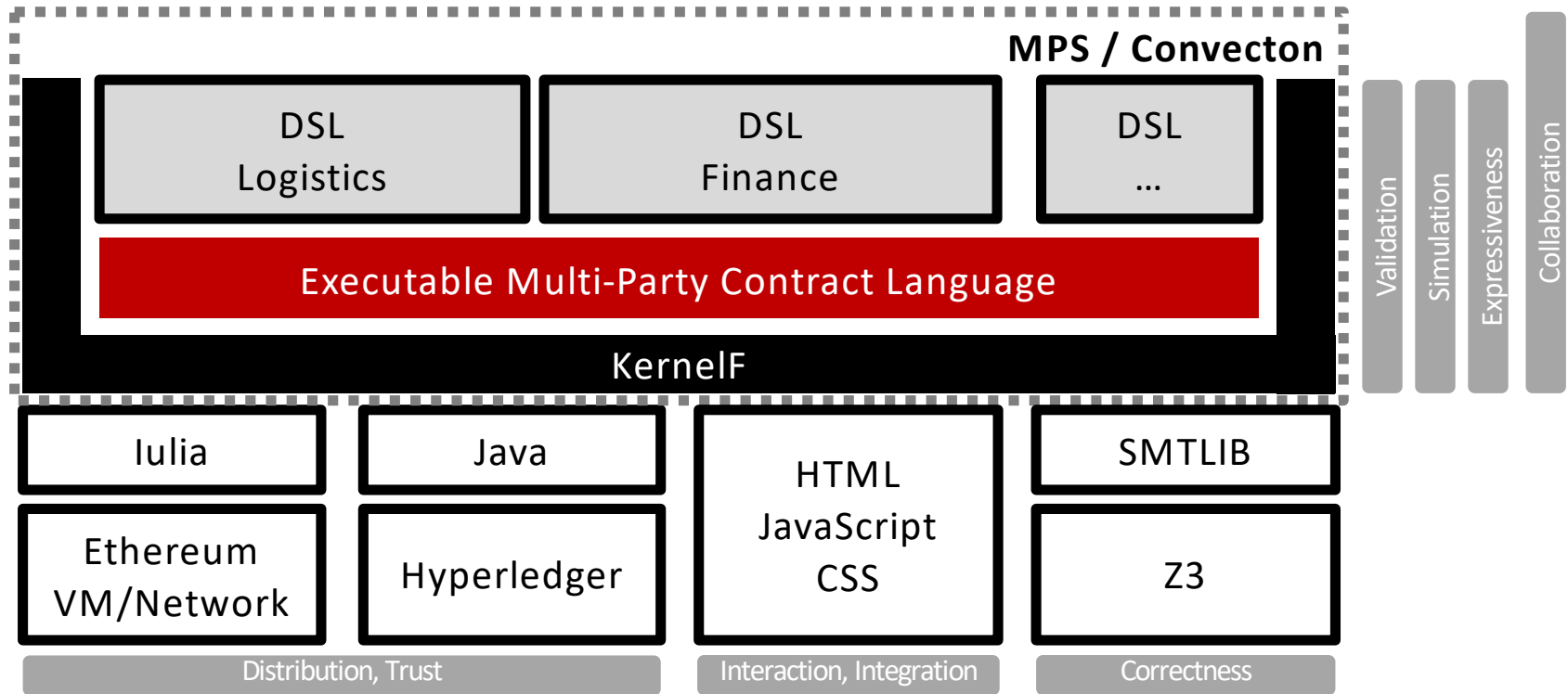
Solution

# An Architecture For Smart Contracts



Generate to verification tools to build more confidence beyond type checking.

# An Architecture For Smart Contracts



# DEMO

# Declarative Description

multi-party-decision Unanimous	
parties: bernd, markus	
dynamic? <input type="checkbox"/>	
procedure: unanimous	time limit: <none>
turnout: <none>	revokable? <input type="checkbox"/>

```
d. [ ] }
```

- decisionTaken
- toList (BaseConcept in jetbrains.mps.lang.core)
- vote
- whoVoted

multi-party-decision Unanimous	
parties: bernd, markus	
dynamic? <input type="checkbox"/>	
procedure: unanimous	time limit: <none>
turnout: all	revokable? <input type="checkbox"/>

```
d. [ ] }
```

- decision
- toList (BaseConcept in jetbrains.mps.lang.core)
- turnoutAchieved
- voteAgainst
- voteFor
- whoVoted
- whoVotedAgainst
- whoVotedFor

## MultiPartyBooleanDecision

A declarative, configurable specification of how a number of parties makes a (Boolean) decision.

# Execution and Test

```
[0] == run(Unanimous) : Unanimous, effects[reads]
    org.ietf3.core.expr.process.plugin.MultipartyBooleanDecisionValue

[1, x] == live($0) : live<Unanimous>
    MultipartyBooleanDecisionValue (snapshot)
    whoVoted -> (collection|0)
    isSealed -> false
    registeredParties -> (collection|2)
                        @[09583503534]
                        @[lfd0g98d09g8sdf]

    decisionTaken -> false

[2] == x.addParty(klaus) : void, effects[modifies]
    MultipartyBooleanDecisionValue (snapshot)
    whoVoted -> (collection|0)
    isSealed -> false
    registeredParties -> (collection|3)
                        @[dsfdslfd0g98d09g8sdf]
                        @[09583503534]
                        @[lfd0g98d09g8sdf]

    decisionTaken -> false
```

```
[3] == x.vote(markus) : void, effects[modifies]
    MultipartyBooleanDecisionValue (snapshot)
    whoVoted -> (collection|1)
                @[09583503534]
    isSealed -> true
    registeredParties -> (collection|3)
                        @[dsfdslfd0g98d09g8sdf]
                        @[09583503534]
                        @[lfd0g98d09g8sdf]

    decisionTaken -> false

[4] == x.vote(klaus) : void, effects[modifies]
    MultipartyBooleanDecisionValue (snapshot)
    whoVoted -> (collection|2)
                @[dsfdslfd0g98d09g8sdf]
                @[09583503534]
    isSealed -> true
    registeredParties -> (collection|3)
                        @[dsfdslfd0g98d09g8sdf]
                        @[09583503534]
                        @[lfd0g98d09g8sdf]

    decisionTaken -> true
```

A MPBD instance maintains the state of a decision process as it evolves over time.

Here, we play with an instance in the interactive REPL.



# Combination with State Machines

More complex contracts are modeled as state machines; events are the API.

```
event openAccess // go to the mode where we allow new guys to request to join
event requestAccess(newGuy: party) // a new guy wants to join the deciders
event terminateAccessRequest(who: party, newGuy: party) // kill a decision procedure
event voteForAccess(voter: party, newGuy: party) // vote for a new guy to become decider
event letsSell // go to the state where we maintain the sell/no-sell decision
event voteForSelling(who: party) // vote for the sale decision
event voteForStopSelling(who: party) // vote against the sale decision
```

Internally, the use BPBDs.

multi-party-decision Sale	
initial parties: bernd, klaus	
dynamic? <input checked="" type="checkbox"/>	sealable? <input type="checkbox"/>
procedure: unanimous	time limit: <none>
turnout: <none>	revokable? <input checked="" type="checkbox"/>

multi-party-decision AccessControl	
initial parties: bernd, klaus	
dynamic? <input checked="" type="checkbox"/>	sealable? <input type="checkbox"/>
procedure: majority	time limit: 20000
turnout: <none>	revokable? <input type="checkbox"/>

```
var sale = run(Sale)
var pendingAccess = box(map<party, AccessControl>())
observable query currentlySelling = sale.decisionTaken
```

# Combination with State Machines II

Here, a transition action creates a new **AccessControl** instance ...

```
on requestAccess(newGuy) [!isDecider/R(newGuy)] : {  
  val acc = run(AccessControl)  
  pendingAccess.update(it.put(newGuy->acc))  
  acc.addParties(sale.registeredParties)  
}
```

The state of that instance is then used in guard condition for the top level SM.

```
on voteForAccess(voter, newGuy) [isPending/R(newGuy) && isDecider/R(voter)] : {  
  val acc = pendingAccess.val[newGuy]  
  acc.vote(voter)  
  if acc.decisionTaken then {  
    sale.addParty(newGuy)  
    pendingAccess.update(it.remove(newGuy))  
  } else none  
}
```

# Preventing Game Theoretical Attacks

Only „valid“ senders can enter this state.

```
state playing [senderIs(players)] {  
  on offerBid(money) : bids := bids.put(sender->money)  
  ...  
}
```

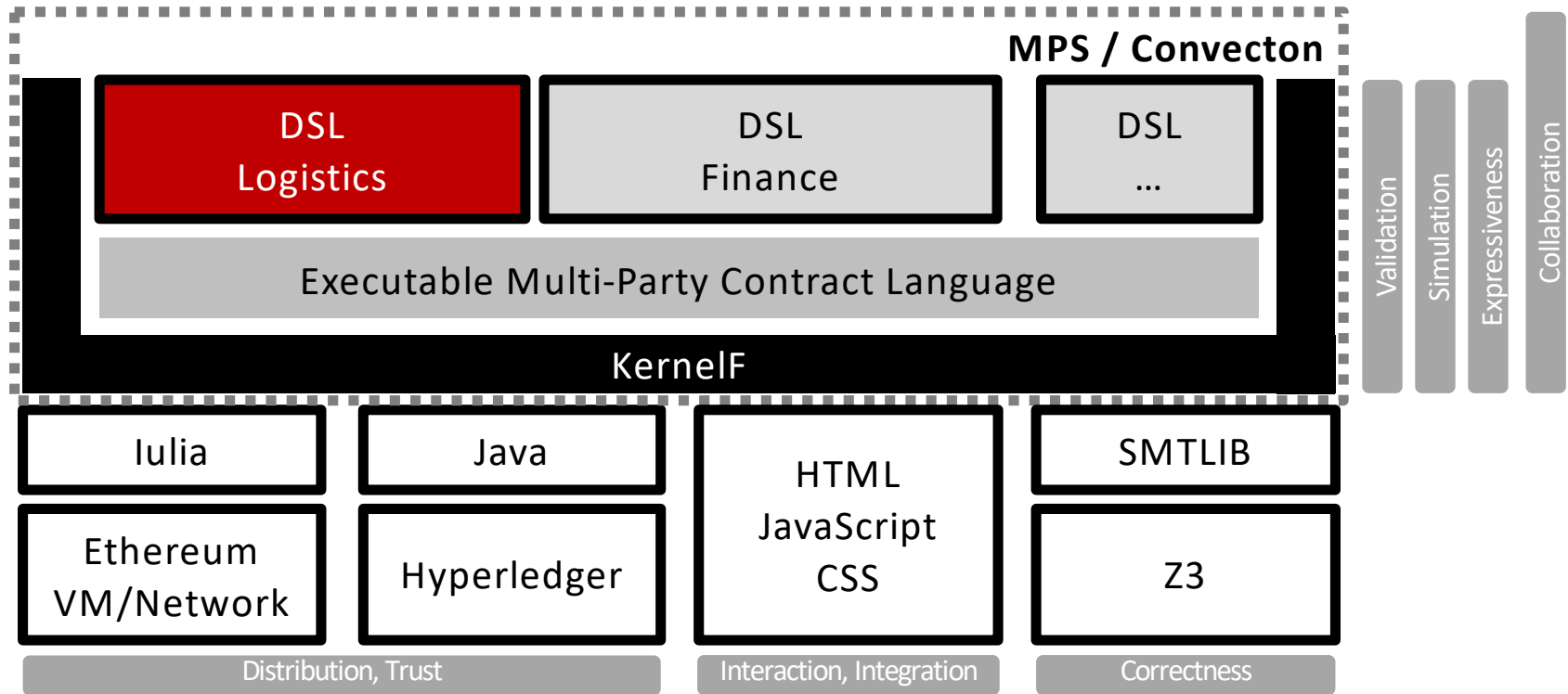
Events can only arrive at limited rate.

```
state requesting [rate(3/1000|commands-only)] {  
  ...  
}
```

States must be entered turn-by-turn.

```
state playing [senderIs(players)] {  
  state bidding [takeTurns(players|ordered|after 1000 remove)] {  
    on offerBid(money) : bids := bids.put(sender->money)  
    if [timeInState > 2000] -> finished  
  }  
  ...  
}
```

# An Architecture For Smart Contracts



# IDEA

# Example: HyperCSL

**SYNGRATO**

Lisp (Clojure) based internal DSL for specification of general commercial contracts.

Inspired by Simon Peyton Jones and Jean-Marc Eber and the POETS group at CPHU and ITU in Denmark.

Uses Ken Adams'  
Categories of Contract



Language as fundamental semantic building blocks. Interpreter and UI in prototype stage.

# SYNGRATO

```
(con :c2
(obs :exercised-timely? :p1)
[(obli :p3
      "lender"
      "borrower"
      (action :payment (event :adv$1000) "Advance loan to borrower.")
      (tw (dt 2014 6 1) (dt 2014 6 2)))
(con :c3
(obs :fulfilled-timely? :p3)
[(obli
   :p4
   "borrower"
   "lender"
   (action :payment (event :pay$550) "Repay first installment.")
   (obs :if (obs :event-occured? :event-of-default)
        (tw (obs :first-time-of :event-of-default))
        (tw (dt 2015 6 1))))
(obli
 :p5
 "borrower"
 "lender"
 (action :payment (event :pay$525) "Repay second installment.")
 (obs :if (obs :event-occured? :event-of-default)
        (tw (obs :first-time-of :event-of-default))
        (tw (dt 2016 6 1))))]]])
```

# Example: HyperCSL

SYNGRATO

## A UI to visualize the interactive execution of CSL contracts.

CSL

### Loan Agreement

c1 - Condition: event-occured? - executed - fulfilled

p0 - Preamble

This loan agreement dated 2014-06-01, by and between Lender Bank Co. and Borrower Corp., will set out the terms under which Lender will extend credit in the principal amount of \$1,000 to Borrower with an un-compounded interest rate of 5% per annum, included in the specified payment structure.

p1 - Discretion - Exercised: true - at: within

Borrower may, by way of notice loan-request within the required time window of: 06/01/2014

p2 - Policy: not-activated

If:

discretion-late?

▪ :p1

then:

agreement-terminated at ?

c2 - Condition: exercised-timely? - p1 - fulfilled

p3 - Obligation - fulfilled: true - at: within

lender must adv\$1000 in favor of borrower within 06/01/2014 - 06/02/2014

c3 - Condition: fulfilled-timely? - p3 - fulfilled

p4 - Obligation - fulfilled: true - at: before

borrower must pay\$550 in favor of lender within 06/04/2015

Input

LOAD CSL

EVALUATE CONTRACT

Clock

06/05/2015

Date

06/05/2015

SET

Log

executed at 06/01/2014



loan-request at 06/01/2014



adv\$1000 at 06/01/2014



pay\$550 at 06/01/2015

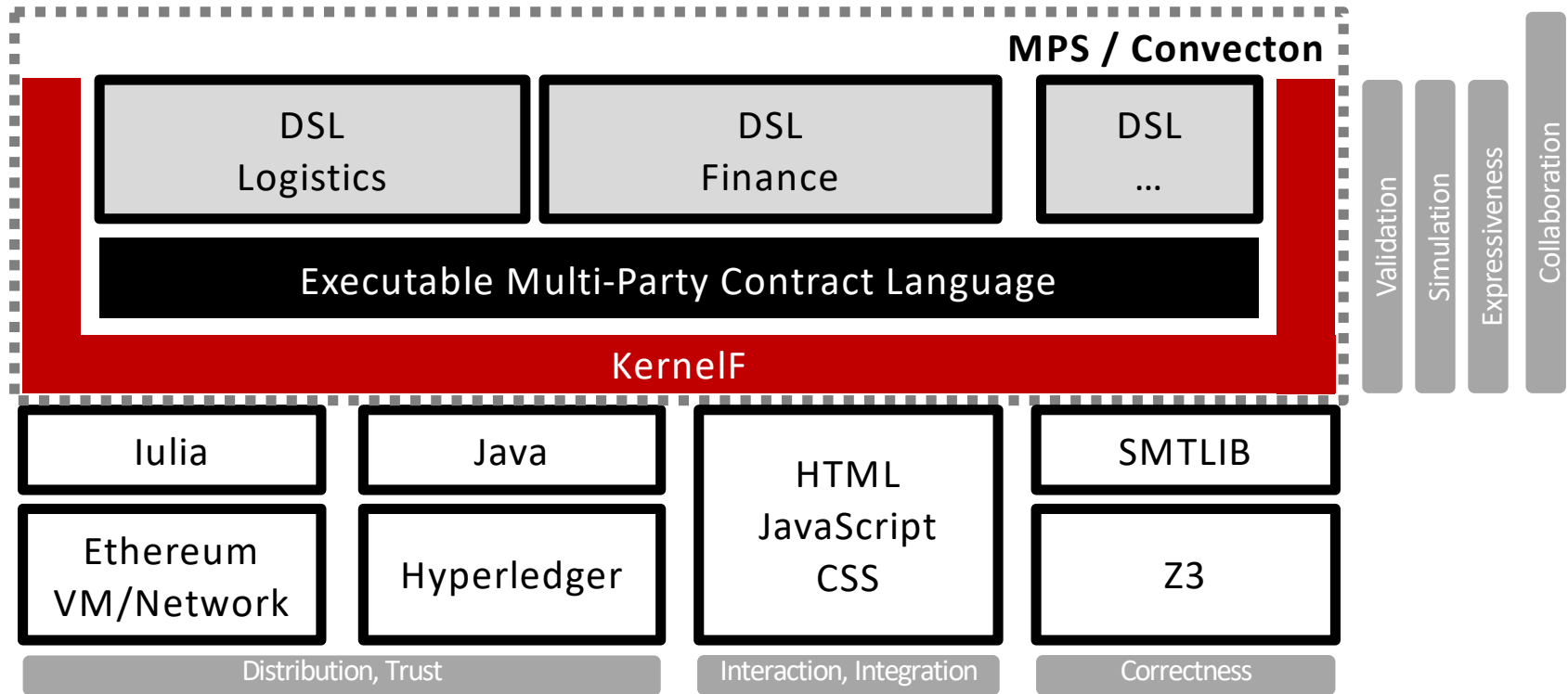


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Tooling



# An Architecture For Smart Contracts



**KernelF is an extensible functional language used at the core of DSLs.**

# DSL Development

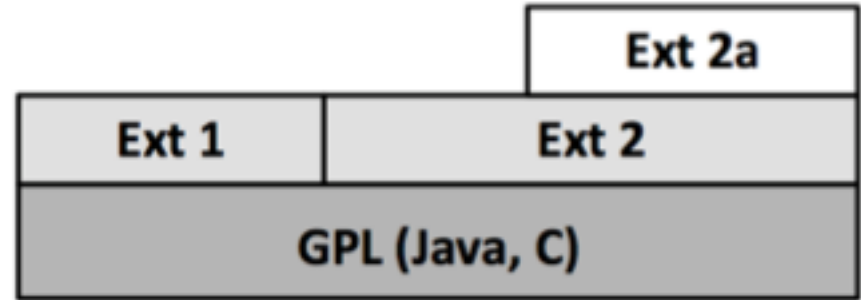
## GPL Extension

Reuse GPL incl. Expressions and TS

Add/Embed DS-extensions

Compatible notational style

Reduce to GPL



## New Language

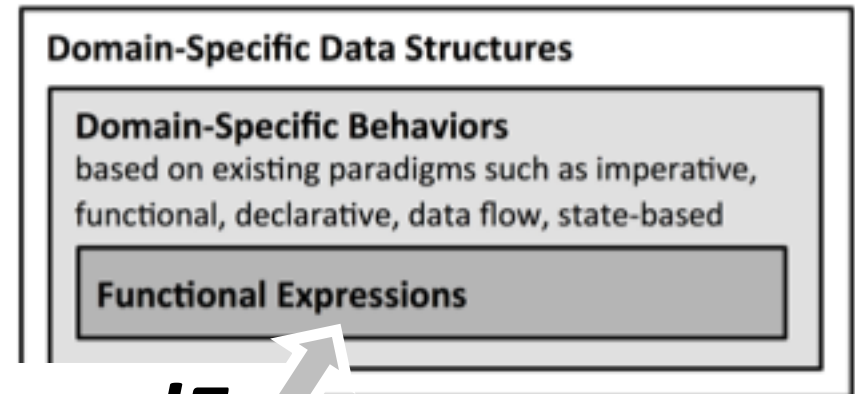
Analyze Domain to find Abstractions

Define suitable, new notations.

Rely on existing behavioral paradigm

Reuse standard expression language

Interpret/Generate to one or more GPLs



**KernelF**

## Formalization

Use existing notation from domain

Clean up and formalize

Generate/Interpret

Often import existing „models“



# Functional Features

Functional, no state at its core.

Purity + Effect Tracking

The usual types, literals and op's

Various Conditionals

Functions and Blocks

No `null`, only `opt<T>`

Error Handling `attempt<T|E-1,... E-n>`

```
try <e> => <s> error <E-1> => <e-1> ... error <E-n> => <e-n>
```

Immutable Collections and higher-order functions

Enums, tuples, records, all immutable

Constraints on types and functions

# Stateful Features

Boxes (like Clojure's `ref`)

Transactional Memory

State Machines

Interactors

Extensible/Embeddable through modular  
language implementation and other means.

# (Meta-) Tooling



Language Workbench

Open Source, by JetBrains

Very Powerful

Used for years by itemis and others

Vast Experience



ZURICH



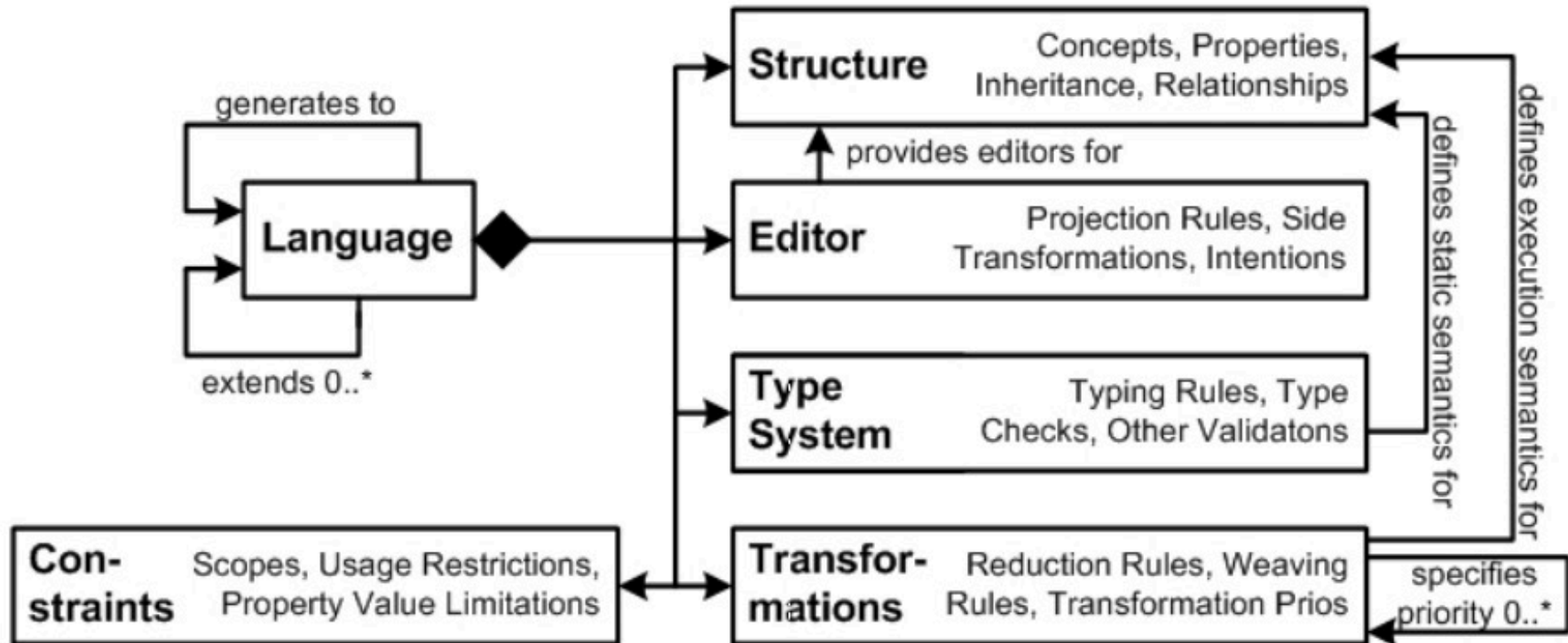
VOLUNTIS



Belastingdienst

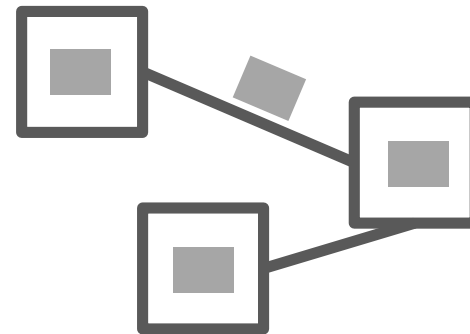


# MPS: Language Toolkit



+ Refactorings, Find Usages, Syntax Coloring, Debugging, ...

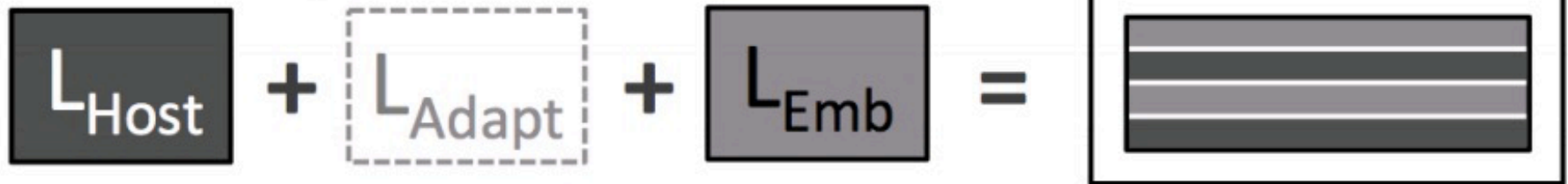
# MPS: Notational Freedom



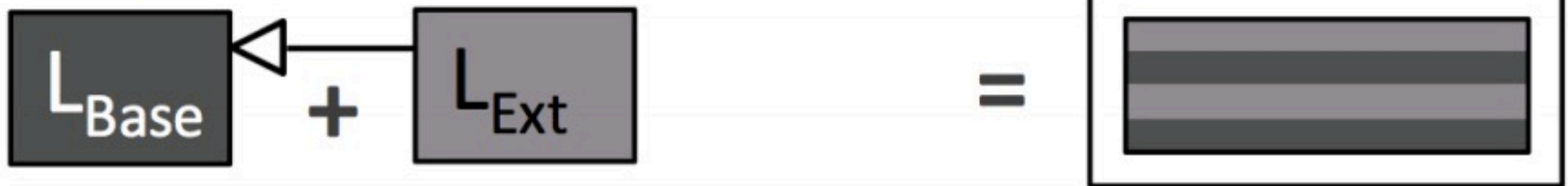
# MPS: Language Composition



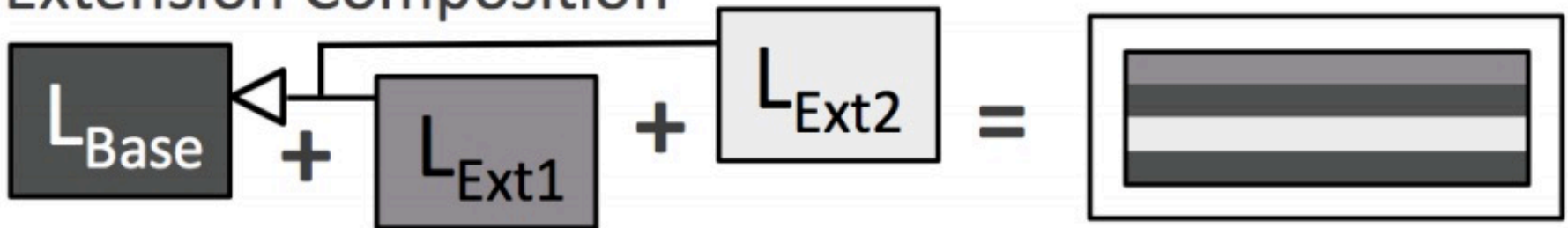
Embedding



Extension



Extension Composition

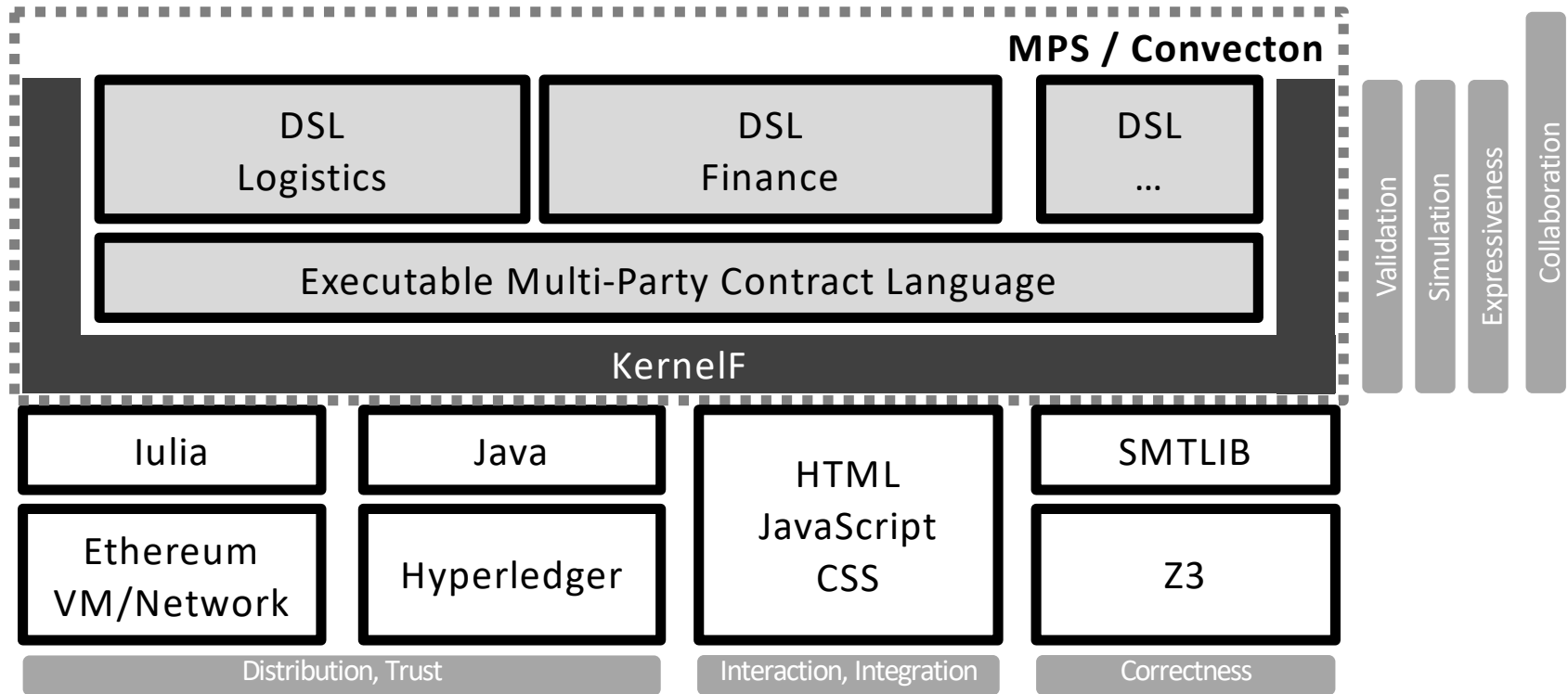




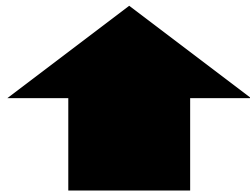
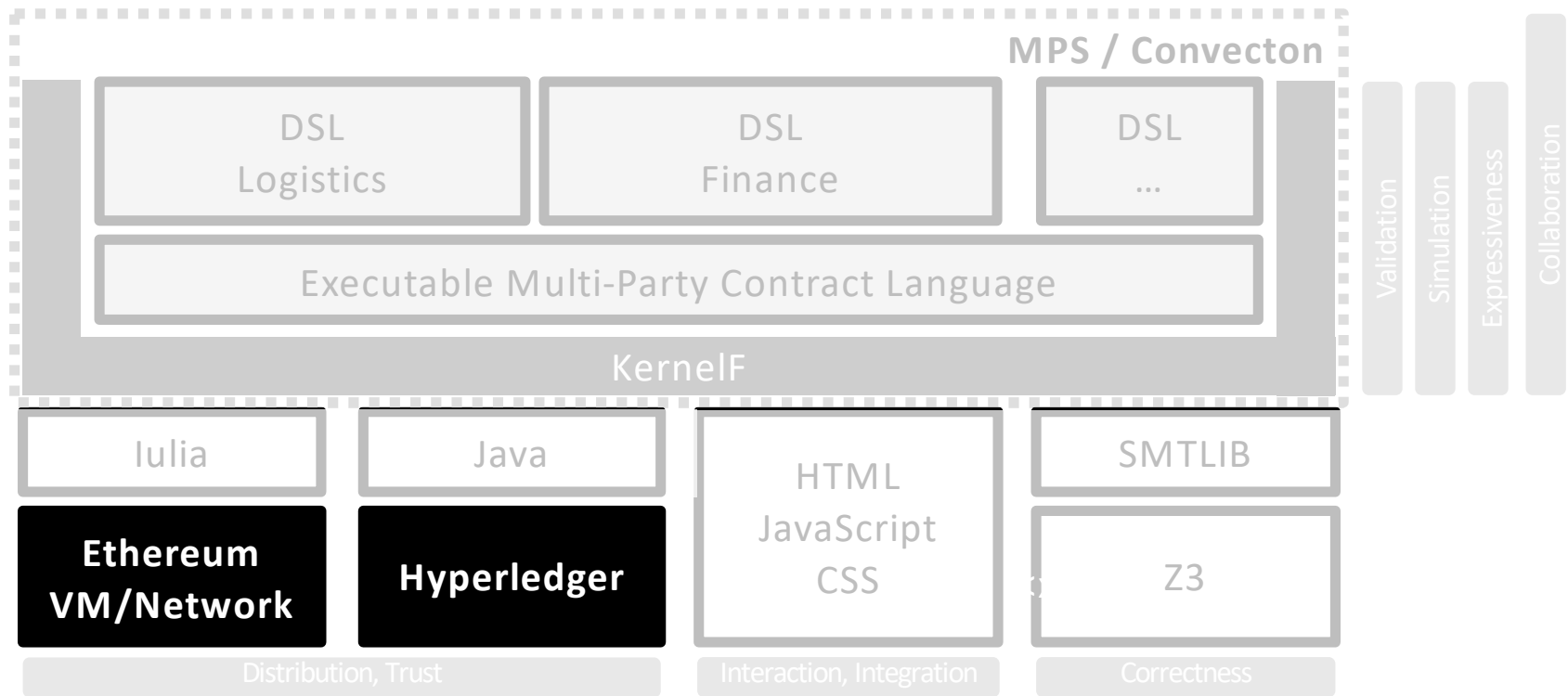
# Verifying Infrastructure



# An Architecture For Smart Contracts



# An Architecture For Smart Contracts



TRUST?

# Verifying Blockchain Infrastructure



## Formal Semantics of the EVM in K

[https://www.ideals.illinois.edu/bitstream/handle/2142/97207/hildenbrandt-saxena-zhu-rodrigues-guth-daian-rosu-2017-tr\\_0818.pdf](https://www.ideals.illinois.edu/bitstream/handle/2142/97207/hildenbrandt-saxena-zhu-rodrigues-guth-daian-rosu-2017-tr_0818.pdf)



## IELE: Register-Based VM for the Blockchain

<https://runtimeverification.com/blog/new-technologies-for-the-blockchain-iele-virtual-machine-and-k-universal-language-framework/>



## ERC20-K: Formal Executable Spec of ERC20

<https://github.com/runtimeverification/erc20-semantics>



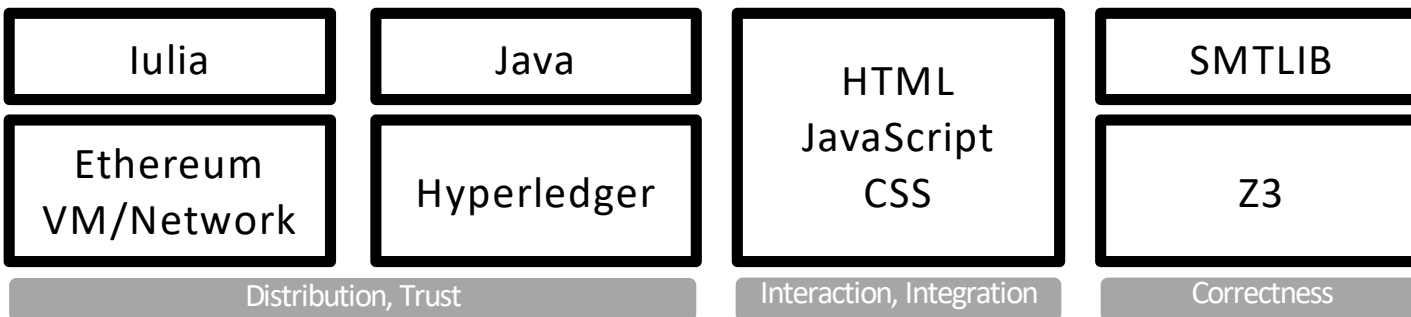
## Formal Verification for Solidity Contracts

<https://forum.ethereum.org/discussion/3779/formal-verification-for-solidity-contracts>

# An Architecture For Smart Contracts



Did I program/specify the right behaviors?



Will the infrastructure execute the behaviors faithfully?

# An Architecture For Smart Contracts

## Validation



Ensure that the program does the correct things, wrt. to the requirements.

Did I program/specify the right behaviors?

## Verification

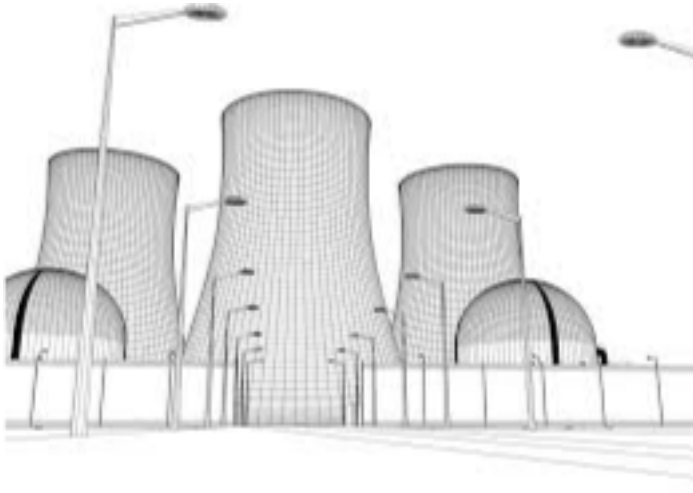
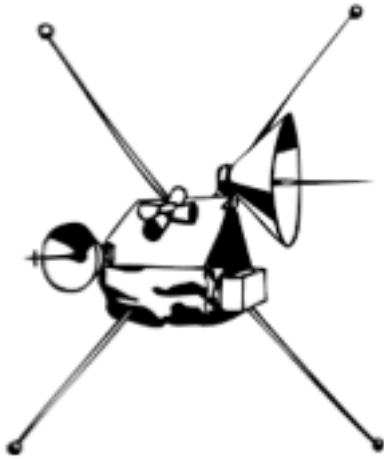


Ensure that the program performs correctly the things the program text tells it to do.

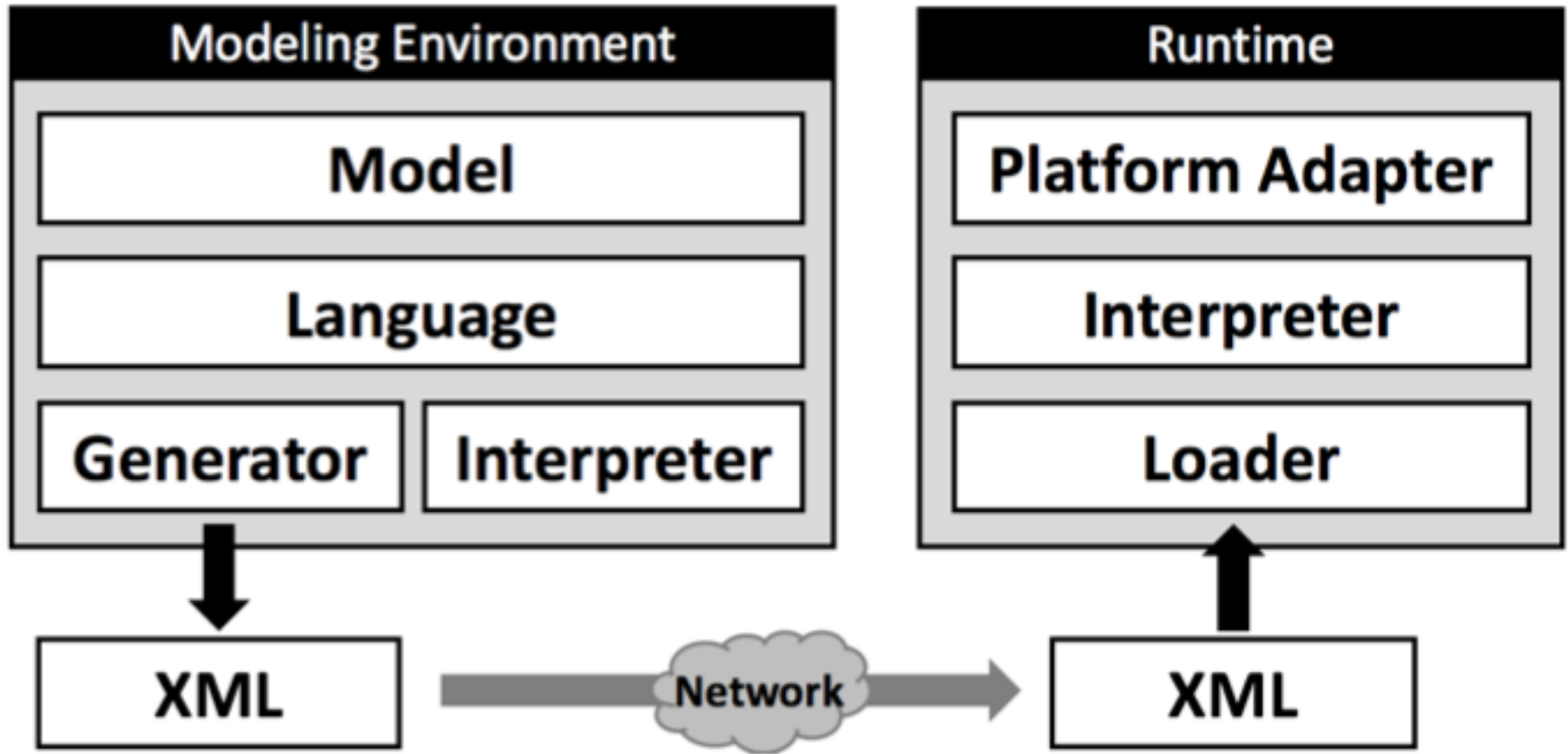
Will the infrastructure execute the behaviors faithfully?



# Not the first community to realize ... 😊



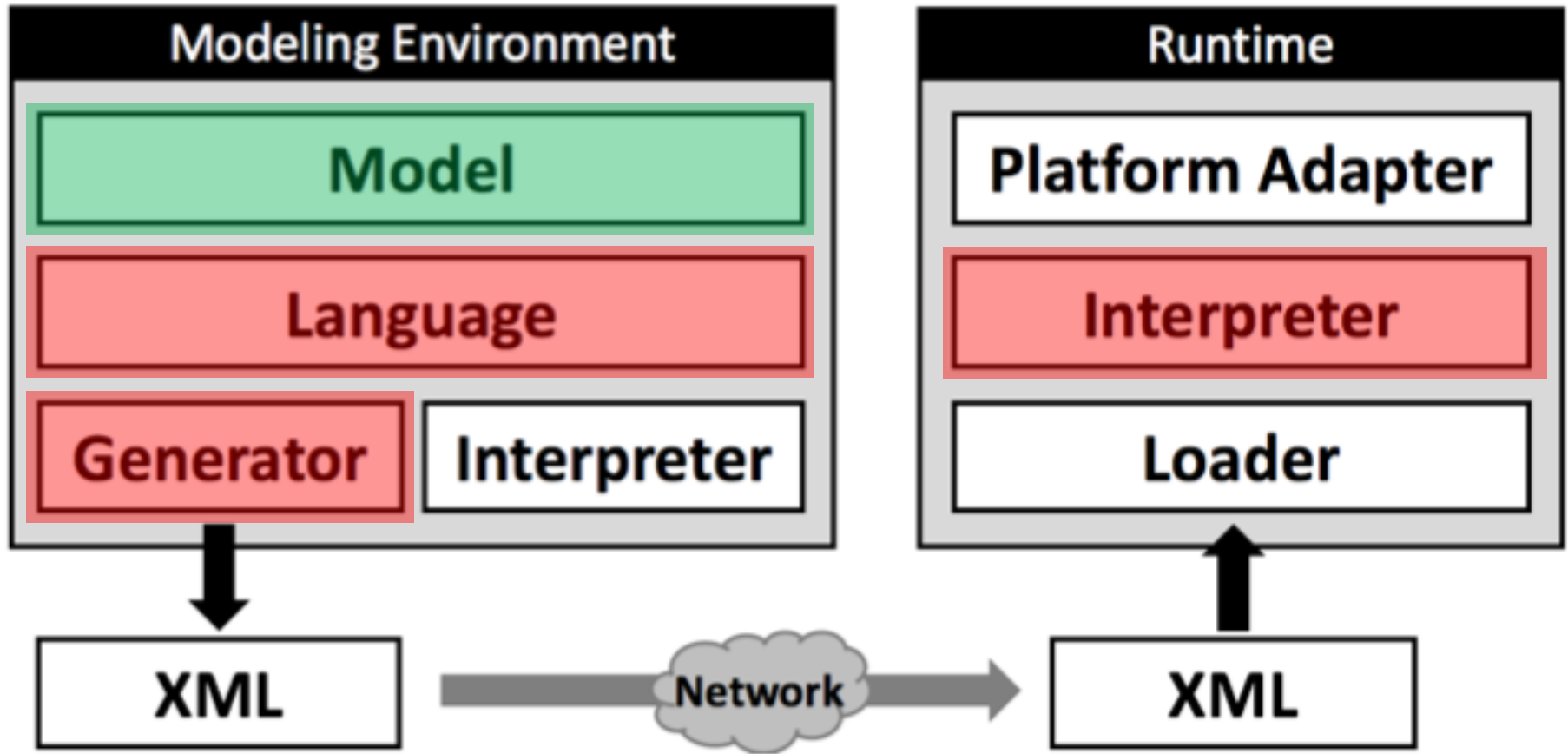
# Sidebar: System Architecture



**What good is all the abstraction if we cannot trust the translation to the implementation?**

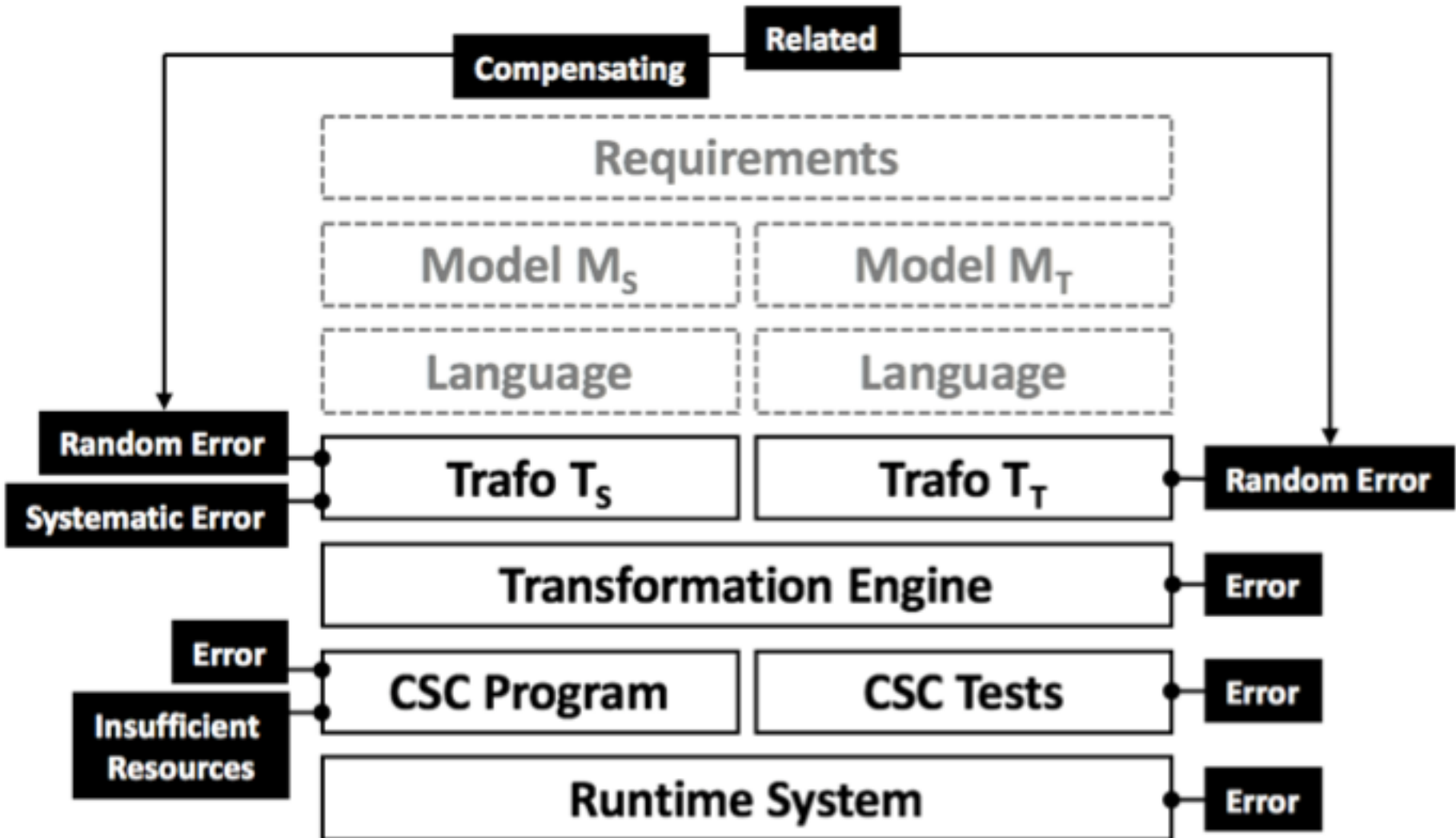


# Sidebar: System Architecture

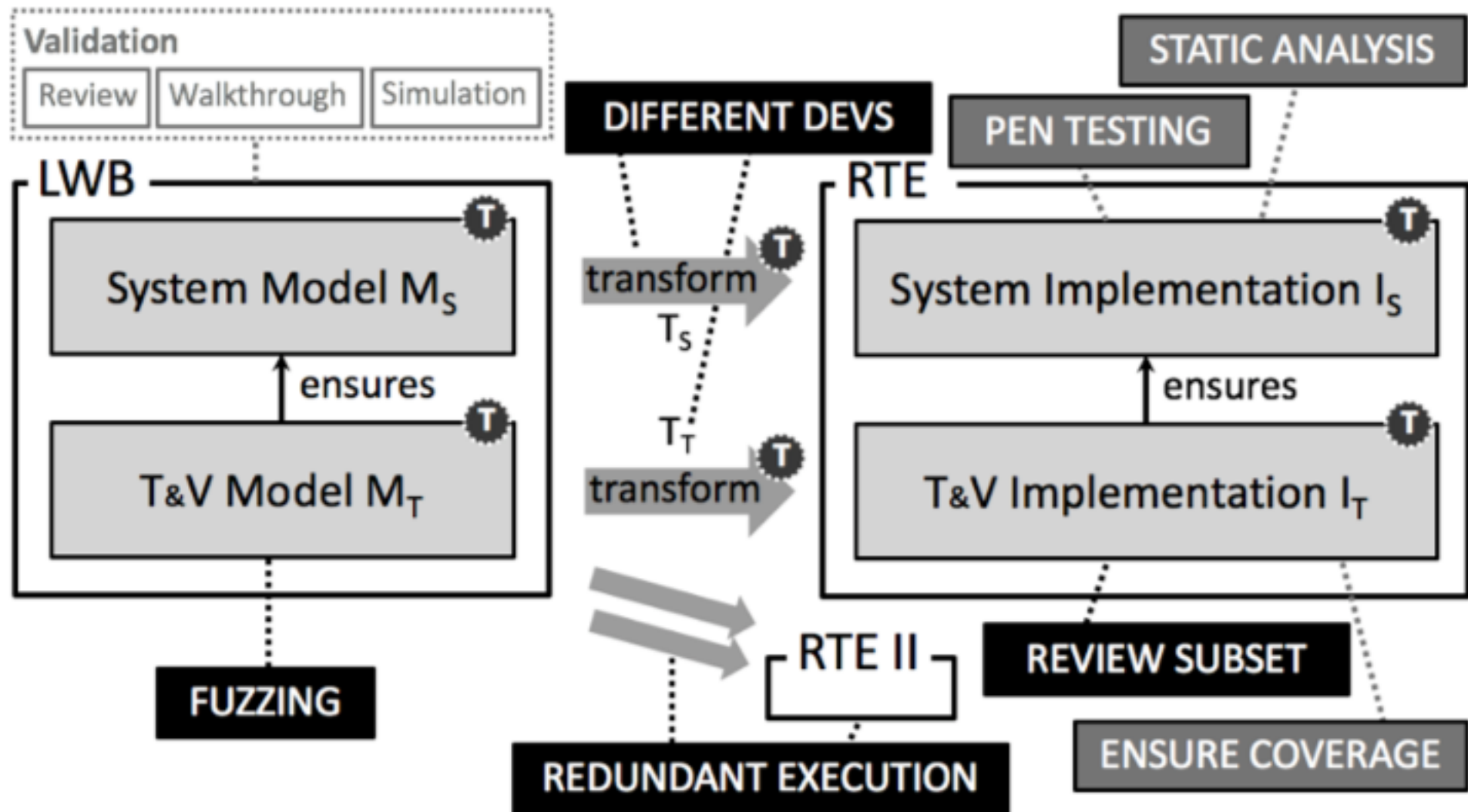


**Tools** may introduce *additional systematic errors* if faulty. Safety **standards** require reliable mitigation of such errors.

# Risk Analysis



# Mitigations – Safe Architecture





# Wrap Up

# Further Reading



## Mutable State in KernelF

[https://medium.com/@markusvoelter/  
dealing-with-mutable-state-in-kernelf-e0fdec8a489b](https://medium.com/@markusvoelter/dealing-with-mutable-state-in-kernelf-e0fdec8a489b)



## A Smart Contract Development Stack

[https://languageengineering.io/  
a-smart-contract-development-stack-54533a3a503a](https://languageengineering.io/a-smart-contract-development-stack-54533a3a503a)



## A Smart Contract Development Stack, Pt. 2

[https://languageengineering.io/a-smart-contract-development-  
stack-part-ii-game-theoretical-aspects-ca7a9d2e548d](https://languageengineering.io/a-smart-contract-development-stack-part-ii-game-theoretical-aspects-ca7a9d2e548d)



## KernelF Reference

<http://voelter.de/data/pub/kernelf-reference.pdf>



## DSLs in Safety-Critical Development

<http://voelter.de/data/pub/MPS-in-Safety-1.0.pdf>

**Contracts must be functionally correct**  
in order for stakeholders to trust them.

**We need better languages**  
to describe contracts in a meaningful way



**Integration of verification tools**  
can be an important step to assure correctness

**Simulation, Experimentation and Test**  
should be available in an interactive, local environment

**Deployment to Blockchain is non-func,**  
it provides guarantees beyond functionality

**Other deployments are useful,**  
that provide other trade-offs (secure↓, fast↑)